

Modular Arithmetic

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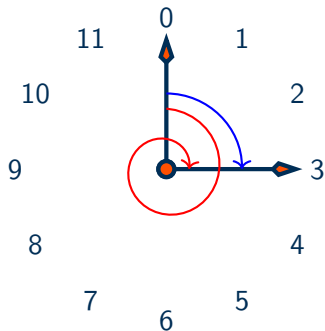
Modular Equivalence

Definition 1

Two numbers a and b are **equivalent modulo n** if $|b - a|$ is a multiple of n .

Consider the 12 hour clock on the left:

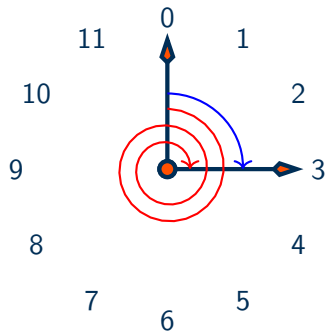
- $3 \equiv 15 \pmod{12}$ because $15 - 3 = 12$



Modular Equivalence

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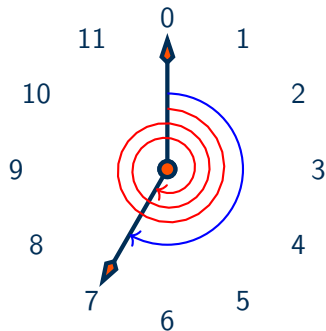
- $3 \equiv 15 \pmod{12}$ because $15 - 3 = 12$
- $3 \equiv 27 \pmod{12}$ because $27 - 3 = 2 \cdot 12$



Modular Equivalence

Definition 1

Two numbers a and b are **equivalent modulo n** if $|b - a|$ is a multiple of n .



Consider the 12 hour clock on the left:

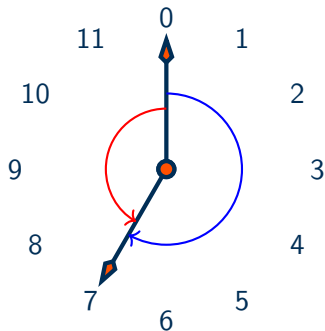
- $3 \equiv 15 \pmod{12}$ because $15 - 3 = 12$
- $3 \equiv 27 \pmod{12}$ because $27 - 3 = 2 \cdot 12$
- $7 \equiv 31 \pmod{12}$ because $31 - 7 = 2 \cdot 12$



Modular Equivalence

Definition 1

Two numbers a and b are **equivalent modulo n** if $|b - a|$ is a multiple of n .



Consider the 12 hour clock on the left:

- $3 \equiv 15 \pmod{12}$ because $15 - 3 = 12$
- $3 \equiv 27 \pmod{12}$ because $27 - 3 = 2 \cdot 12$
- $7 \equiv 31 \pmod{12}$ because $31 - 7 = 2 \cdot 12$
- $7 \equiv -5 \pmod{12}$ because $7 - (-5) = 12$



Modular Arithmetic

Theorem 2

Given $n \in \mathbb{N}$ and $a, b, c, d \in \mathbb{Z}$, if $a \equiv b \pmod{n}$ and $c \equiv d \pmod{N}$, then

$$a + c \equiv b + d \pmod{n}$$

$$a - c \equiv b - d \pmod{n}$$

$$ac \equiv bd \pmod{n}$$



Arithmetic Modulo 7

$+_7$	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

\times_7	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1



Arithmetic Modulo 10

$+_{10}$	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	0
2	2	3	4	5	6	7	8	9	0	1
3	3	4	5	6	7	8	9	0	1	2
4	4	5	6	7	8	9	0	1	2	3
5	5	6	7	8	9	0	1	2	3	4
6	6	7	8	9	0	1	2	3	4	5
7	7	8	9	0	1	2	3	4	5	6
8	8	9	0	1	2	3	4	5	6	7
9	9	0	1	2	3	4	5	6	7	8

\times_{10}	0	1	2	3	4	5	6	7	8	9
0	0	0	0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6	7	8	9
2	0	2	4	6	8	0	2	4	6	8
3	0	3	6	9	2	5	8	1	4	7
4	0	4	8	2	6	0	4	8	2	6
5	0	5	0	5	0	5	0	5	0	5
6	0	6	2	8	4	0	6	2	8	4
7	0	7	4	1	8	5	2	9	6	3
8	0	8	6	4	2	0	8	6	4	2
9	0	9	8	7	6	5	4	3	2	1



Proof Modular Arithmetic is Well-Defined

Proof of Theorem 2.

Let $n \in \mathbb{N}$ and $a, b, c, d \in \mathbb{Z}$, let $a \equiv b \pmod{n}$ and $c \equiv d \pmod{n}$. By definition, $a - b = q_1n$ and $c - d = q_2n$ so that

$$(a + c) - (b + d) = (q_1 + q_2)n;$$

thus

$$a + c \equiv b + d \pmod{n}.$$



Proof Modular Arithmetic is Well-Defined

Proof of Theorem 2.

Let $n \in \mathbb{N}$ and $a, b, c, d \in \mathbb{Z}$, let $a \equiv b \pmod{n}$ and $c \equiv d \pmod{n}$. By definition, $a - b = q_1n$ and $c - d = q_2n$ so that

$$(ac) - (bd) = (a - b)c + (c - d)b = (q_1c + q_2b)n;$$

thus

$$ac \equiv bd \pmod{n}.$$



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Identities, Inverses, and Zero Divisors

Definition 3

The **additive identity** is 0 because $a + 0 = 0 + a = a$ for any a . And, we say that a and b are **additive inverses** if $a + b = 0$.



Identities, Inverses, and Zero Divisors

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The **additive identity** is 0 because $a + 0 = 0 + a = a$ for any a . And, we say that a and b are **additive inverses** if $a + b = 0$.

Definition 4

The **multiplicative identity** is 1 because $a \cdot 1 = 1 \cdot a = a$ for any a . And, we say that a and b are **multiplicative inverses** if $a \cdot b = 1$.



Identities, Inverses, and Zero Divisors

Definition 3

The **additive identity** is 0 because $a + 0 = 0 + a = a$ for any a . And, we say that a and b are **additive inverses** if $a + b = 0$.

Definition 4

The **multiplicative identity** is 1 because $a \cdot 1 = 1 \cdot a = a$ for any a . And, we say that a and b are **multiplicative inverses** if $a \cdot b = 1$.

Definition 5

We say that $a \neq 0$ is a **zero divisor** if there exists $b \neq 0$ with $a \cdot b = 0$.



Arithmetic Modulo 7

$+_7$	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

\times_7	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1



Arithmetic Modulo 7

$+7$	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

$\times 7$	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1



Arithmetic Modulo 10

+10	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	0
2	2	3	4	5	6	7	8	9	0	1
3	3	4	5	6	7	8	9	0	1	2
4	4	5	6	7	8	9	0	1	2	3
5	5	6	7	8	9	0	1	2	3	4
6	6	7	8	9	0	1	2	3	4	5
7	7	8	9	0	1	2	3	4	5	6
8	8	9	0	1	2	3	4	5	6	7
9	9	0	1	2	3	4	5	6	7	8

×10	0	1	2	3	4	5	6	7	8	9
0	0	0	0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6	7	8	9
2	0	2	4	6	8	0	2	4	6	8
3	0	3	6	9	2	5	8	1	4	7
4	0	4	8	2	6	0	4	8	2	6
5	0	5	0	5	0	5	0	5	0	5
6	0	6	2	8	4	0	6	2	8	4
7	0	7	4	1	8	5	2	9	6	3
8	0	8	6	4	2	0	8	6	4	2
9	0	9	8	7	6	5	4	3	2	1



Arithmetic Modulo 10

+10	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	0
2	2	3	4	5	6	7	8	9	0	1
3	3	4	5	6	7	8	9	0	1	2
4	4	5	6	7	8	9	0	1	2	3
5	5	6	7	8	9	0	1	2	3	4
6	6	7	8	9	0	1	2	3	4	5
7	7	8	9	0	1	2	3	4	5	6
8	8	9	0	1	2	3	4	5	6	7
9	9	0	1	2	3	4	5	6	7	8

×10	0	1	2	3	4	5	6	7	8	9
0	0	0	0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6	7	8	9
2	0	2	4	6	8	0	2	4	6	8
3	0	3	6	9	2	5	8	1	4	7
4	0	4	8	2	6	0	4	8	2	6
5	0	5	0	5	0	5	0	5	0	5
6	0	6	2	8	4	0	6	2	8	4
7	0	7	4	1	8	5	2	9	6	3
8	0	8	6	4	2	0	8	6	4	2
9	0	9	8	7	6	5	4	3	2	1



Bezout and Inverses

Theorem 6 (Bezout's Theorem)

Given $a, b \in \mathbb{Z}$, $d = (a, b)$ if and only if $d = ax + by$ is the least positive linear combination of a and b .



Bezout and Inverses

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Given $a, b \in \mathbb{Z}$, $d = (a, b)$ if and only if $d = ax + by$ is the least positive linear combination of a and b .

Theorem 7 (Modular Inverses)

Given $n \in \mathbb{N}$, $a \in \mathbb{Z}$, and $d = (a, n)$, a has a multiplicative inverse modulo n if $d = 1$, otherwise a is a zero divisor.



Proof of 7

Proof.

If $(a, n) = 1$, then by theorem 6, $ax + ny = 1$ and $ax \equiv 1 \pmod{n}$, i.e. $x \equiv a^{-1} \pmod{n}$.



Proof of 7

Proof.

If $(a, n) = 1$, then by theorem 6, $ax + ny = 1$ and $ax \equiv 1 \pmod{n}$, i.e. $x \equiv a^{-1} \pmod{n}$.

Suppose $(a, n) = d \neq 1$, then by theorem 6, $ax + ny = d$. Let $q = n/d < n$, so that

$$aqx \equiv qd \equiv 0 \pmod{n},$$

i.e. a is a zero divisor. □



Comments on Some Equivalencies

Theorem 8

If r_1, r_2, \dots, r_n is a **complete set of residues** modulo n , $(a, n) = 1$, and $b \in \mathbb{Z}$ then

$$ar_1 + b, ar_2 + b, \dots, ar_n + b$$

is also a complete set of residues.



Comments on Some Equivalencies

Theorem 9

If $a, b, k, m \in \mathbb{Z}$ with $k, m > 0$ and $a \equiv b \pmod{m}$, then

$$a^k \equiv b^k \pmod{m}.$$



Comments on Some Equivalencies

Theorem 9

If $a, b, k, m \in \mathbb{Z}$ with $k, m > 0$ and $a \equiv b \pmod{m}$, then

$$a^k \equiv b^k \pmod{m}.$$

Theorem 10

If $a, b, m_1, m_2, \dots, m_k \in \mathbb{Z}$ with $m_i > 0$ and $a \equiv b \pmod{m_i}$ for all i , then

$$a \equiv b \pmod{[m_1, m_2, \dots, m_k]}.$$



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Solving Linear Congruences

Theorem 11 (Solving Single Linear Congruences)

Let $n \in \mathbb{N}$, and $a, b, c \in \mathbb{Z}$, then

$$ax + b \equiv c \pmod{n},$$

has solution $x \equiv a^{-1}(c - b) \pmod{n}$ provided $a^{-1} \pmod{n}$ exists.



Congruences Modulo 7

Example 12

Find x such that $5x + 2 \equiv 1 \pmod{7}$.

Congruences Modulo 7

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Find x such that $5x + 2 \equiv 1 \pmod{7}$.

Modulo 7, $5^{-1} \equiv 3$, so we get

$$x \equiv 3(1 - 2) \equiv 4 \pmod{7}.$$

Congruences Modulo 7

Example 12

Find x such that $5x + 2 \equiv 1 \pmod{7}$.

Modulo 7, $5^{-1} \equiv 3$, so we get

$$x \equiv 3(1 - 2) \equiv 4 \pmod{7}.$$

Check:

$$\begin{aligned} 5 \cdot 4 + 2 &= 22 \\ &\equiv 1 \pmod{7}. \end{aligned}$$

Congruences Modulo 10

Example 13

Find x such that $3x + 2 \equiv 5 \pmod{10}$.



Congruences Modulo 10

Example 13

Find x such that $3x + 2 \equiv 5 \pmod{10}$.

Modulo 10, $3^{-1} \equiv 7$, so we get

$$x \equiv 7(5 - 2) \equiv 1 \pmod{10}.$$



Congruences Modulo 10

Example 13

Find x such that $3x + 2 \equiv 5 \pmod{10}$.

Modulo 10, $3^{-1} \equiv 7$, so we get

$$x \equiv 7(5 - 2) \equiv 1 \pmod{10}.$$

Check:

$$3 \cdot 1 + 2 \equiv 5 \pmod{10}.$$



Congruences Modulo 10

Example 14

Find x such that $2x + 3 \equiv 5 \pmod{10}$.



Congruences Modulo 10

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Find x such that $2x + 3 \equiv 5 \pmod{10}$.

Modulo 10, 2 doesn't have a multiplicative inverse, but we can easily see that $x = 1$ still works.



Congruences Modulo 10

Example 14

Find x such that $2x + 3 \equiv 5 \pmod{10}$.

Modulo 10, 2 doesn't have a multiplicative inverse, but we can easily see that $x = 1$ still works.

Example 15

Find x such that $2x \equiv 5 \pmod{10}$.



Congruences Modulo 10

Example 14

Find x such that $2x + 3 \equiv 5 \pmod{10}$.

Modulo 10, 2 doesn't have a multiplicative inverse, but we can easily see that $x = 1$ still works.

Example 15

Find x such that $2x \equiv 5 \pmod{10}$.

Modulo 10, 2 doesn't have a multiplicative inverse and in this case there is no solution.



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Affine Cipher: $C \equiv m(M) + s \pmod{26}$

Algorithm

- 1 Key: m, s # a multiplier and a shift
- 2 Message Character: M
- 3 Replace M with $C \equiv m \cdot M + s \pmod{26}$
- 4 Repeat step 3 for each character



Arithmetic Modulo 26

\times_{26}	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
2	0	2	4	6	8	10	12	14	16	18	20	22	24	0	2	4	6	8	10	12	14	16	18	20	22	24
3	0	3	6	9	12	15	18	21	24	1	4	7	10	13	16	19	22	25	2	5	8	11	14	17	20	23
4	0	4	8	12	16	20	24	2	6	10	14	18	22	0	4	8	12	16	20	24	2	6	10	14	18	22
5	0	5	10	15	20	25	4	9	14	19	24	3	8	13	18	23	2	7	12	17	22	1	6	11	16	21
6	0	6	12	18	24	4	10	16	22	2	8	14	20	0	6	12	18	24	4	10	16	22	2	8	14	20
7	0	7	14	21	2	9	16	23	4	11	18	25	6	13	20	1	8	15	22	3	10	17	24	5	12	19
8	0	8	16	24	6	14	22	4	12	20	2	10	18	0	8	16	24	6	14	22	4	12	20	2	10	18
9	0	9	18	1	10	19	2	11	20	3	12	21	4	13	22	5	14	23	6	15	24	7	16	25	8	17
10	0	10	20	4	14	24	8	18	2	12	22	6	16	0	10	20	4	14	24	8	18	2	12	22	6	16
11	0	11	22	7	18	3	14	25	10	21	6	17	2	13	24	9	20	5	16	1	12	23	8	19	4	15
12	0	12	24	10	22	8	20	6	18	4	16	2	14	0	12	24	10	22	8	20	6	18	4	16	2	14
13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13
14	0	14	2	16	4	18	6	20	8	22	10	24	12	0	14	2	16	4	18	6	20	8	22	10	24	12
15	0	15	4	19	8	23	12	1	16	5	20	9	24	13	2	17	6	21	10	25	14	3	18	7	22	11
16	0	16	6	22	12	2	18	8	24	14	4	20	10	0	16	6	22	12	2	18	8	24	14	4	20	10
17	0	17	8	25	16	7	24	15	6	23	14	5	22	13	4	21	12	3	20	11	2	19	10	1	18	9
18	0	18	10	2	20	12	4	22	14	6	24	16	8	0	18	10	2	20	12	4	22	14	6	24	16	8
19	0	19	12	5	24	17	10	3	22	15	8	1	20	13	6	25	18	11	4	23	16	9	2	21	14	7
20	0	20	14	8	2	22	16	10	4	24	18	12	6	0	20	14	8	2	22	16	10	4	24	18	12	6
21	0	21	16	11	6	1	22	17	12	7	2	23	18	13	8	3	24	19	14	9	4	25	20	15	10	5
22	0	22	18	14	10	6	2	24	20	16	12	8	4	0	22	18	14	10	6	2	24	20	16	12	8	4
23	0	23	20	17	14	11	8	5	2	25	22	19	16	13	10	7	4	1	24	21	18	15	12	9	6	3
24	0	24	22	20	18	16	14	12	10	8	6	4	2	0	24	22	20	18	16	14	12	10	8	6	4	2
25	0	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1



Arithmetic Modulo 26

\times_{26}	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
2	0	2	4	6	8	10	12	14	16	18	20	22	24	0	2	4	6	8	10	12	14	16	18	20	22	24
3	0	3	6	9	12	15	18	21	24	1	4	7	10	13	16	19	22	25	2	5	8	11	14	17	20	23
4	0	4	8	12	16	20	24	2	6	10	14	18	22	0	4	8	12	16	20	24	2	6	10	14	18	22
5	0	5	10	15	20	25	4	9	14	19	24	3	8	13	18	23	2	7	12	17	22	1	6	11	16	21
6	0	6	12	18	24	4	10	16	22	2	8	14	20	0	6	12	18	24	4	10	16	22	2	8	14	20
7	0	7	14	21	2	9	16	23	4	11	18	25	6	13	20	1	8	15	22	3	10	17	24	5	12	19
8	0	8	16	24	6	14	22	4	12	20	2	10	18	0	8	16	24	6	14	22	4	12	20	2	10	18
9	0	9	18	1	10	19	2	11	20	3	12	21	4	13	22	5	14	23	6	15	24	7	16	25	8	17
10	0	10	20	4	14	24	8	18	2	12	22	6	16	0	10	20	4	14	24	8	18	2	12	22	6	16
11	0	11	22	7	18	3	14	25	10	21	6	17	2	13	24	9	20	5	16	1	12	23	8	19	4	15
12	0	12	24	10	22	8	20	6	18	4	16	2	14	0	12	24	10	22	8	20	6	18	4	16	2	14
13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13
14	0	14	2	16	4	18	6	20	8	22	10	24	12	0	14	2	16	4	18	6	20	8	22	10	24	12
15	0	15	4	19	8	23	12	1	16	5	20	9	24	13	2	17	6	21	10	25	14	3	18	7	22	11
16	0	16	6	22	12	2	18	8	24	14	4	20	10	0	16	6	22	12	2	18	8	24	14	4	20	10
17	0	17	8	25	16	7	24	15	6	23	14	5	22	13	4	21	12	3	20	11	2	19	10	1	18	9
18	0	18	10	2	20	12	4	22	14	6	24	16	8	0	18	10	2	20	12	4	22	14	6	24	16	8
19	0	19	12	5	24	17	10	3	22	15	8	1	20	13	6	25	18	11	4	23	16	9	2	21	14	7
20	0	20	14	8	2	22	16	10	4	24	18	12	6	0	20	14	8	2	22	16	10	4	24	18	12	6
21	0	21	16	11	6	1	22	17	12	7	2	23	18	13	8	3	24	19	14	9	4	25	20	15	10	5
22	0	22	18	14	10	6	2	24	20	16	12	8	4	0	22	18	14	10	6	2	24	20	16	12	8	4
23	0	23	20	17	14	11	8	5	2	25	22	19	16	13	10	7	4	1	24	21	18	15	12	9	6	3
24	0	24	22	20	18	16	14	12	10	8	6	4	2	0	24	22	20	18	16	14	12	10	8	6	4	2
25	0	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1



Arithmetic Modulo 26

\times_{26}	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
2	0	2	4	6	8	10	12	14	16	18	20	22	24	0	2	4	6	8	10	12	14	16	18	20	22	24
3	0	3	6	9	12	15	18	21	24	1	4	7	10	13	16	19	22	25	2	5	8	11	14	17	20	23
4	0	4	8	12	16	20	24	2	6	10	14	18	22	0	4	8	12	16	20	24	2	6	10	14	18	22
5	0	5	10	15	20	25	4	9	14	19	24	3	8	13	18	23	2	7	12	17	22	1	6	11	16	21
6	0	6	12	18	24	4	10	16	22	2	8	14	20	0	6	12	18	24	4	10	16	22	2	8	14	20
7	0	7	14	21	2	9	16	23	4	11	18	25	6	13	20	1	8	15	22	3	10	17	24	5	12	19
8	0	8	16	24	6	14	22	4	12	20	2	10	18	0	8	16	24	6	14	22	4	12	20	2	10	18
9	0	9	18	1	10	19	2	11	20	3	12	21	4	13	22	5	14	23	6	15	24	7	16	25	8	17
10	0	10	20	4	14	24	8	18	2	12	22	6	16	0	10	20	4	14	24	8	18	2	12	22	6	16
11	0	11	22	7	18	3	14	25	10	21	6	17	2	13	24	9	20	5	16	1	12	23	8	19	4	15
12	0	12	24	10	22	8	20	6	18	4	16	2	14	0	12	24	10	22	8	20	6	18	4	16	2	14
13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13	0	13
14	0	14	2	16	4	18	6	20	8	22	10	24	12	0	14	2	16	4	18	6	20	8	22	10	24	12
15	0	15	4	19	8	23	12	1	16	5	20	9	24	13	2	17	6	21	10	25	14	3	18	7	22	11
16	0	16	6	22	12	2	18	8	24	14	4	20	10	0	16	6	22	12	2	18	8	24	14	4	20	10
17	0	17	8	25	16	7	24	15	6	23	14	5	22	13	4	21	12	3	20	11	2	19	10	1	18	9
18	0	18	10	2	20	12	4	22	14	6	24	16	8	0	18	10	2	20	12	4	22	14	6	24	16	8
19	0	19	12	5	24	17	10	3	22	15	8	1	20	13	6	25	18	11	4	23	16	9	2	21	14	7
20	0	20	14	8	2	22	16	10	4	24	18	12	6	0	20	14	8	2	22	16	10	4	24	18	12	6
21	0	21	16	11	6	1	22	17	12	7	2	23	18	13	8	3	24	19	14	9	4	25	20	15	10	5
22	0	22	18	14	10	6	2	24	20	16	12	8	4	0	22	18	14	10	6	2	24	20	16	12	8	4
23	0	23	20	17	14	11	8	5	2	25	22	19	16	13	10	7	4	1	24	21	18	15	12	9	6	3
24	0	24	22	20	18	16	14	12	10	8	6	4	2	0	24	22	20	18	16	14	12	10	8	6	4	2
25	0	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1



Affine Cipher: $C \equiv m(M) + s \pmod{26}$

Algorithm

- 1 Key: m, s # a multiplier and a shift
- 2 Message Character: M
- 3 Replace M with $C \equiv m \cdot M + s \pmod{26}$
- 4 Repeat step 3 for each character

Key Alphabet with Multiplier of $m=7$ and $s=4$

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
plain	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
CIPHER	E	L	S	Z	G	N	U	B	I	P	W	D	K	R	Y	F	M	T	A	H	O	V	C	J	Q	X



Affine Cipher: $C \equiv m(M) + s \pmod{26}$

Algorithm

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plain	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
CIPHER	E	L	S	Z	G	N	U	B	I	P	W	D	K	R	Y	F	M	T	A	H	O	V	C	J	Q	X

$$7 \cdot \text{"d"} + 4 = 7 \cdot 3 + 4 = 25 \equiv 25 \pmod{26} \quad (\text{or } C = \text{"Z"})$$



Affine Cipher: $C \equiv m(M) + s \pmod{26}$

Algorithm

- 1 Key: m, s # a multiplier and a shift
- 2 Message Character: M
- 3 Replace M with $C \equiv m \cdot M + s \pmod{26}$
- 4 Repeat step 3 for each character

Key Alphabet with Multiplier of $m=7$ and $s=4$

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
plain	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
CIPHER	E	L	S	Z	G	N	U	B	I	P	W	D	K	R	Y	F	M	T	A	H	O	V	C	J	Q	X

$$7 \cdot \text{"m"} + 4 = 7 \cdot 12 + 4 = 88 \equiv 10 \pmod{26} \text{ (or } C=\text{"K"})$$



Affine Cipher: $C \equiv m(M) + s \pmod{26}$

Algorithm

- 1 Key: m, s # a multiplier and a shift
- 2 Message Character: M
- 3 Replace M with $C \equiv m \cdot M + s \pmod{26}$
- 4 Repeat step 3 for each character

Key Alphabet with Multiplier of $m=7$ and $s=4$

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
plain	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
CIPHER	E	L	S	Z	G	N	U	B	I	P	W	D	K	R	Y	F	M	T	A	H	O	V	C	J	Q	X

$$7 \cdot \text{"p"} + 4 = 7 \cdot 15 + 4 = 109 \equiv 5 \pmod{26} \text{ (or } C = \text{"F"})$$



Inverting an Affine Cipher

Invert the Cipher

Given

$$C = m \cdot M + s$$



Inverting an Affine Cipher

Invert the Cipher

Given

$$C = m \cdot M + s$$

we can solve for C:

$$m \cdot M + s = C \Rightarrow m \cdot M = C - s$$



Inverting an Affine Cipher

Invert the Cipher

Given

$$C = m \cdot M + s$$

we can solve for C:

$$\begin{aligned} m \cdot M + s = C &\Rightarrow m \cdot M = C - s \\ &\Rightarrow M = m^{-1}(C - s) \end{aligned}$$



Inverting an Affine Cipher

Invert the Cipher

Given

$$C = m \cdot M + s$$

we can solve for C:

$$m \cdot M + s = C \Rightarrow m \cdot M = C - s$$

$$\Rightarrow M = m^{-1}(C - s)$$

$$\Rightarrow M = m^{-1}C - m^{-1}s$$



Inverting an Affine Cipher

Invert the Cipher

Given

$$C = m \cdot M + s$$

we can solve for C:

$$\begin{aligned} m \cdot M + s = C &\Rightarrow m \cdot M = C - s \\ &\Rightarrow M = m^{-1}(C - s) \\ &\Rightarrow M = m^{-1}C - m^{-1}s \end{aligned}$$

so we have two expressions we can use to find M.



Deciphering: $M \equiv m^{-1}C - m^{-1}s \pmod{26}$

Algorithm using $M \equiv m^{-1}C - m^{-1}s \pmod{26}$

- 1 Key = m, s # a multiplier and a shift
- 2 Inverse Key = $m^{-1}, -s$ # mult. and add. inverses
- 3 Simplify Key: $m^{-1}, -m^{-1}s$
- 4 Cipher Character = C
- 5 Replace C with $M \equiv m^{-1}C + (-m^{-1}s) \pmod{26}$
- 6 Repeat step 5 for each cipher character



Sample Deciphering

Deciphering with an Algorithm

Cipher Text: "IDGET RGDZY RUEUY RGVGT HYCTG AHDGC IHBEF IUQYO
UGHZI THQER ZLGAI ZGAHB GFIUD IWGAI H"

- 1 Key: $m=7, s=4$
- 2 Inverse Key: $m^{-1} = 15, -s = -4 \equiv 22$
- 3 Simplified Key: $m^{-1} = 15, -m^{-1}s = 15 \cdot 22 \equiv 18$
- 4 $C = "I" = 8$
- 5 $M \equiv 15 \cdot 8 + 18 \equiv 8 \pmod{26}, M = "i"$
- 6 plain message = "i"
- 7 Next $C = "D" = 3$



Sample Deciphering

Deciphering with an Algorithm

Cipher Text: "IDGET RGDZY RUEUY RGVGT HYCTG AHDGC IHBEF IUQYO
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Sample Deciphering

Deciphering with an Algorithm

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- ⑤ $M \equiv 15 \cdot 8 + 18 \equiv 8 \pmod{26}, M = "i"$
- ⑥ plain message = "i"
- ⑦ Next $C = "D" = 3$



Sample Deciphering

Deciphering with an Algorithm

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- 3 Simplified Key: $m^{-1} = 15, -m^{-1}s = 15 \cdot 22 \equiv 18$
- 4 $C = "D" = 3$
- 5 $M \equiv 15 \cdot 3 + 18 \equiv 11 \pmod{26}, M = "I"$
- 6 plain message = "il"
- 7 Next $C = "G" = 6$



Sample Deciphering

Deciphering with an Algorithm

Cipher Text: "IDGET RGDZY RUEUY RGVGT HYCTG AHDGC IHBEF IUQYO
UGHZI THQER ZLGAI ZGAHB GFIUD IWGAI H"

- 1 Key: $m=7, s=4$
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Sample Deciphering

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- 4 $C = "G" = 6$
- 5 $M \equiv 15 \cdot 6 + 18 \equiv 4 \pmod{26}, M = "e"$
- 6 plain message = "ile"
- 7 Next $C = "E" = 4$



Sample Deciphering

Deciphering with an Algorithm

Cipher Text: "IDGET RYZDY RUEUY RYVGT HYCTG AHDGC IHBEF IUQYO
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- 4 $C = "E" = 4$
- 5 $M \equiv 15 \cdot 4 + 18 \equiv 0 \pmod{26}, M = "a"$
- 6 plain message = "ilea"
- 7 Next $C = "T" = 19$



Sample Deciphering

Deciphering with an Algorithm

Cipher Text: "IDGET RGDZY RUEUY RGVGT HYCTG AHDGC IHBEF IUQYO
UGHZI THQER ZLGAI ZGAHB GFIUD IWGAI H"

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- ④ $C = "E" = 4$
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- ⑥ plain message = "ilea"
- ⑦ Next $C = "T" = 19$



Sample Deciphering

Deciphering with an Algorithm

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- 6 plain message = "ilea"
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Sample Deciphering

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Sample Deciphering

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- 3 Simplified Key: $m^{-1} = 15, -m^{-1}s = 15 \cdot 22 \equiv 18$
- 4 $C = "T" = 19$
- 5 $M \equiv 15 \cdot 19 + 18 \equiv 17 \pmod{26}, M = "r"$
- 6 plain message = "ilear"
- 7 Next $C = "R" = 17$



Sample Deciphering

Deciphering with an Algorithm

Cipher Text: "IDGET RGVGT RUEUY RGVGT HYCTG AHDGC IHBEF IUQYO
UGHZI THQER ZLGAI ZGAHB GFIUD IWGAI H"

- 1 Key: $m=7, s=4$
- 2 Inverse Key: $m^{-1} = 15, -s = -4 \equiv 22$
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- ⑥ plain message = "ilear"
- ⑦ Next $C = "R" = 17$



Sample Deciphering

Deciphering with an Algorithm

Cipher Text: "IDGET RGZDY RUEUY RGVGT HYCTG AHDGC IHBEF IUQYO
UGHZI THQER ZLGAI ZGAHB GFIUD IWGAI H"

- ① Key: $m=7, s=4$
- ② Inverse Key: $m^{-1} = 15, -s = -4 \equiv 22$
- ③ Simplified Key: $m^{-1} = 15, -m^{-1}s = 15 \cdot 22 \equiv 18$
- ④ $C = "H" = 7$
- ⑤ $M \equiv 15 \cdot 7 + 18 \equiv 19 \pmod{26}, M = "t"$
- ⑥ plain message = "ilearnedlongagonevertowrestlewithapigyouget-dirtyandbesidesthepiglikesit"
- ⑦ final message = "I learned long ago, never to wrestle with a pig. You get dirty, and besides, the pig likes it."

Sample Deciphering

Deciphering with an Algorithm

Cipher Text: "IDGET RGZDY RUEUY RGVGT HYCTG AHDGC IHBEF IUQYO
UGHZI THQER ZLGAI ZGAHB GFIUD IWGAI H"

- ① Key: $m=7, s=4$
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Chinese Remainder Theorem

Theorem 16 (Chinese Remainder Theorem)

Given the set of linear equations

$$x \equiv a_1 \pmod{n_1}$$

$$x \equiv a_2 \pmod{n_2}$$

$$x \equiv a_3 \pmod{n_3}$$

$$\vdots$$

$$x \equiv a_r \pmod{n_r},$$

provided the n_i are relatively prime, there exists a unique solution modulo $M = n_1 n_2 n_3 \cdots n_r$.

C.R.T. Proof

Proof.

Let $M = n_1 n_2 n_3 \cdots n_r$ and $M_i = M/n_i$. Since the n_i are relatively prime, $(M_i, n_i) = 1$ and we can calculate $M_i^{-1} \pmod{n_i}$.



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Let $M = n_1 n_2 n_3 \cdots n_r$ and $M_i = M/n_i$. Since the n_i are relatively prime, $(M_i, n_i) = 1$ and we can calculate $M_i^{-1} \pmod{n_i}$. Then

$$x \equiv \sum a_i M_i M_i^{-1} \pmod{M},$$

is the desired solution.



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$$x \equiv \sum a_i M_i M_i^{-1} \pmod{M},$$

is the desired solution. Note that $M_j \equiv 0 \pmod{n_i}$ if $i \neq j$ and $M_i M_i^{-1} \equiv 1 \pmod{n_i}$ so that $x \equiv a_i \pmod{n_i}$. □



C.R.T. Proof

Proof.

Let $M = n_1 n_2 n_3 \cdots n_r$ and $M_i = M/n_i$. Since the n_i are relatively prime, $(M_i, n_i) = 1$ and we can calculate $M_i^{-1} \pmod{n_i}$. Then

$$x \equiv \sum a_i M_i M_i^{-1} \pmod{M},$$

is the desired solution. Note that $M_j \equiv 0 \pmod{n_i}$ if $i \neq j$ and $M_i M_i^{-1} \equiv 1 \pmod{n_i}$ so that $x \equiv a_i \pmod{n_i}$. Finally if there existed two solutions x_1 and x_2 then theorem 10 ensures they are equal modulo M . \square



C.R.T. Example 1

$$x \equiv 3 \pmod{5}$$

$$x \equiv 1 \pmod{6}$$

$$x \equiv 2 \pmod{7}$$



C.R.T. Example 1

$$x \equiv 3 \pmod{5}$$

$$x \equiv 1 \pmod{6}$$

$$x \equiv 2 \pmod{7}$$

$$M = 210$$



C.R.T. Example 1

$$x \equiv 3 \pmod{5}$$

$$x \equiv 1 \pmod{6}$$

$$x \equiv 2 \pmod{7}$$

$$M = 210$$

$$M_1 = 42$$



C.R.T. Example 1

$$x \equiv 3 \pmod{5}$$

$$x \equiv 1 \pmod{6}$$

$$x \equiv 2 \pmod{7}$$

$$M = 210$$

$$M_1 = 42$$

$$M_2 = 35$$



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$$x \equiv 3 \pmod{5}$$

$$x \equiv 1 \pmod{6}$$

$$x \equiv 2 \pmod{7}$$

$$M = 210$$

$$M_1 = 42$$

$$M_2 = 35$$

$$M_3 = 30$$



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$$M_1^{-1} \equiv 3 \pmod{5}$$



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$$M_3^{-1} \equiv 4 \pmod{7}$$

$$x \equiv \sum a_i M_i M_i^{-1} \pmod{M}$$



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$$M_3^{-1} \equiv 4 \pmod{7}$$

$$\begin{aligned} x &\equiv \sum a_i M_i M_i^{-1} \pmod{M} \\ &\equiv 3 \cdot 42 \cdot 3 + 1 \cdot 35 \cdot 5 + 2 \cdot 30 \cdot 4 \pmod{M} \end{aligned}$$



C.R.T. Example 1

$$x \equiv 3 \pmod{5}$$

$$x \equiv 1 \pmod{6}$$

$$x \equiv 2 \pmod{7}$$

$$M = 210$$

$$M_1 = 42$$

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$$\begin{aligned} x &\equiv \sum a_i M_i M_i^{-1} \pmod{M} \\ &\equiv 3 \cdot 42 \cdot 3 + 1 \cdot 35 \cdot 5 + 2 \cdot 30 \cdot 4 \pmod{M} \\ &\equiv 378 + 175 + 240 \pmod{M} \end{aligned}$$



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$$M = 210$$

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$$x \equiv 3 \pmod{5}$$

$$x \equiv 1 \pmod{6}$$

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$$M = 210$$

$$M_1 = 42$$

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$$x \equiv 3 \pmod{5}$$

$$x \equiv 1 \pmod{6}$$

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$$M = 210$$

$$M_1 = 42$$

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$$\begin{aligned} x &\equiv \sum a_i M_i M_i^{-1} \pmod{M} \\ &\equiv 3 \cdot 42 \cdot 3 + 1 \cdot 35 \cdot 5 + 2 \cdot 30 \cdot 4 \pmod{M} \\ &\equiv 378 + 175 + 240 \pmod{M} \\ &\equiv 169 + -35 + 30 \pmod{M} \\ &\equiv 163 \pmod{M} \end{aligned}$$

Check:

$$163 \equiv 3 \pmod{5}$$

$$163 \equiv 1 \pmod{6}$$

$$163 \equiv 2 \pmod{7}$$



C.R.T. Example 2

$$x \equiv 3 \pmod{6}$$

$$x \equiv 7 \pmod{10}$$



C.R.T. Example 2

$$x \equiv 3 \pmod{6}$$

$$x \equiv 7 \pmod{10}$$

$$x \equiv 1 \pmod{2}$$

$$x \equiv 0 \pmod{3}$$

$$x \equiv 1 \pmod{2}$$

$$x \equiv 2 \pmod{5}$$



C.R.T. Example 2

$$x \equiv 3 \pmod{6}$$

$$x \equiv 7 \pmod{10}$$

$$x \equiv 1 \pmod{2}$$

$$x \equiv 0 \pmod{3}$$

$$x \equiv 1 \pmod{2}$$

$$x \equiv 2 \pmod{5}$$

$$M = 30$$

$$M_1 = 15$$

$$M_2 = 10$$

$$M_3 = 6$$



C.R.T. Example 2

$$x \equiv 3 \pmod{6} \quad M_1^{-1} \equiv 1 \pmod{2}$$

$$x \equiv 7 \pmod{10} \quad M_2^{-1} \equiv 1 \pmod{3}$$

$$x \equiv 1 \pmod{2} \quad M_3^{-1} \equiv 1 \pmod{5}$$

$$x \equiv 0 \pmod{3}$$

$$x \equiv 1 \pmod{2}$$

$$x \equiv 2 \pmod{5}$$

$$M = 30$$

$$M_1 = 15$$

$$M_2 = 10$$

$$M_3 = 6$$



C.R.T. Example 2

$$x \equiv 3 \pmod{6}$$

$$x \equiv 7 \pmod{10}$$

$$x \equiv 1 \pmod{2}$$

$$x \equiv 0 \pmod{3}$$

$$x \equiv 1 \pmod{2}$$

$$x \equiv 2 \pmod{5}$$

$$M = 30$$

$$M_1 = 15$$

$$M_2 = 10$$

$$M_3 = 6$$

$$M_1^{-1} \equiv 1 \pmod{2}$$

$$M_2^{-1} \equiv 1 \pmod{3}$$

$$M_3^{-1} \equiv 1 \pmod{5}$$

$$\begin{aligned} x &\equiv \sum a_i M_i M_i^{-1} \pmod{M} \\ &= 1 \cdot 15 \cdot 1 + 0 \cdot 10 \cdot 1 + 2 \cdot 6 \cdot 1 \pmod{M} \\ &= 27 \pmod{M} \end{aligned}$$



C.R.T. Example 3

$$x \equiv 7 \pmod{15}$$

$$x \equiv 3 \pmod{21}$$

$$x \equiv 4 \pmod{35}$$



C.R.T. Example 3

$$x \equiv 7 \pmod{15}$$

$$x \equiv 3 \pmod{21}$$

$$x \equiv 4 \pmod{35}$$

$$x \equiv 1 \pmod{3}$$

$$x \equiv 2 \pmod{5}$$

$$x \equiv 0 \pmod{3}$$

$$x \equiv 3 \pmod{7}$$

$$x \equiv 4 \pmod{5}$$

$$x \equiv 4 \pmod{7}$$



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Modular Arithmetic

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